

# Recurrent Neural Networks: A Natural Model of Computation Beyond the Turing Limits

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6 October 2012

# Introduction

- ▶ We follow the so-called *mind-computer analogy* approach to cognitive science.
- ▶ We study the computational capabilities of basic models of recurrent neural networks.
- ▶ We show that recurrent neural networks provide a natural model of computation beyond the Turing limits.

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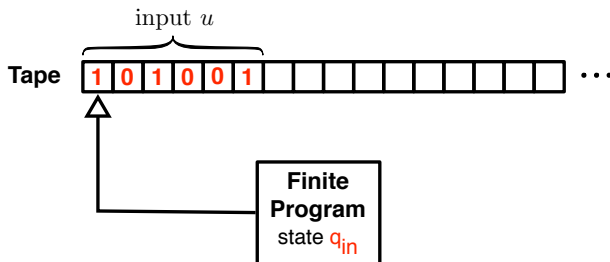
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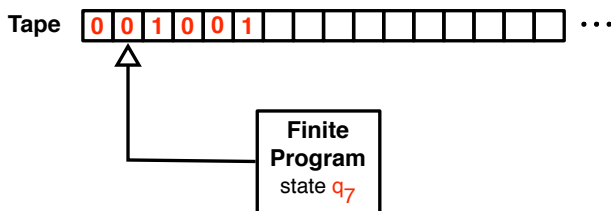
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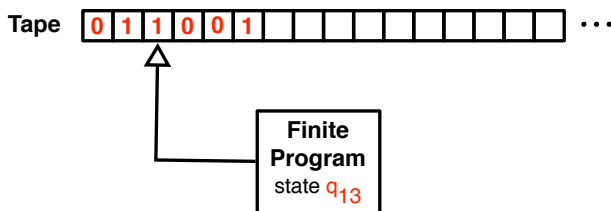
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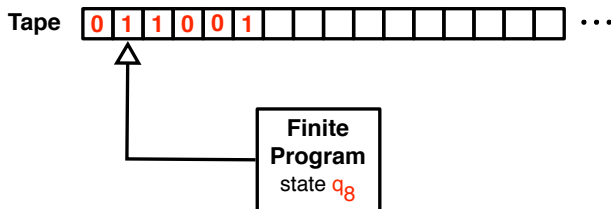
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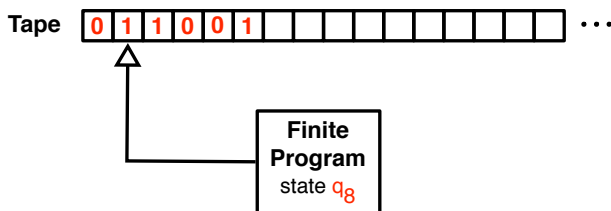


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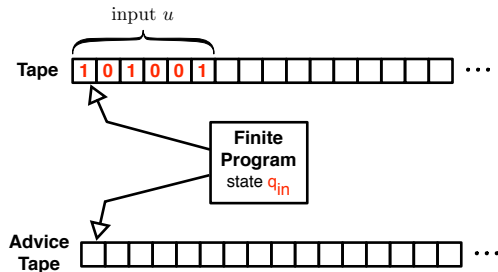
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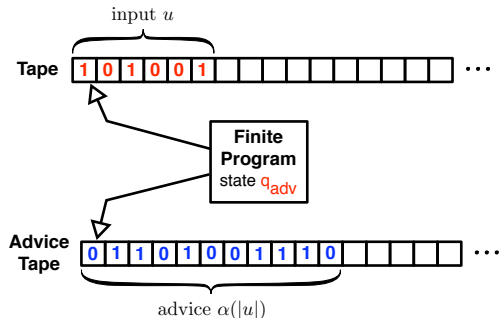
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A Turing machine with advice (TM/A) is a Turing machine provided with an additional advice tape and advice function  $\alpha : \mathbb{N} \longrightarrow \{0, 1\}^*$ .



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# TM vs TM/A

The complexity class **P** (**PTIME**) is the collections of all languages decidable in polynomial time by some TM.

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## Lemma

*TM/poly(A)s are strictly more powerful than TMs (super-Turing) already in polynomial time of computation, i.e.  $\mathbf{P/poly} \supsetneq \mathbf{P}$ .*

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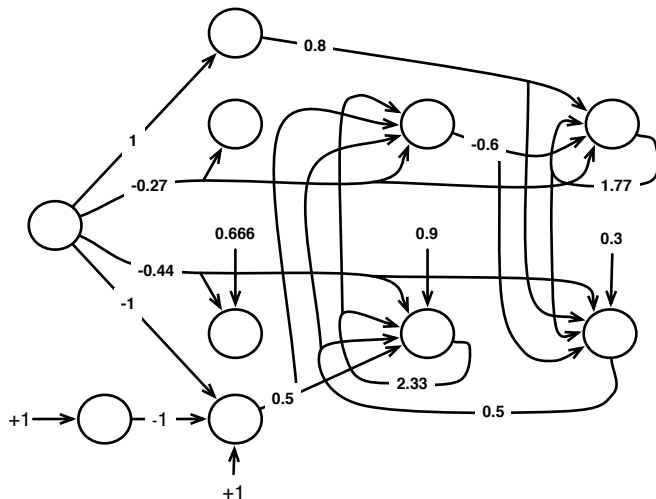
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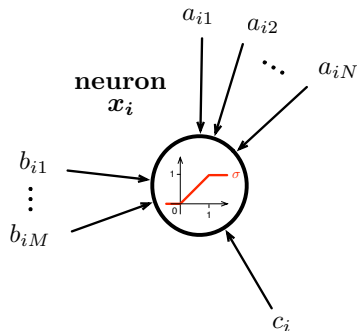
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# Recurrent neural networks



# Dynamics: static synaptic weights



$$x_i(t+1) = \sigma \left( \sum_{j=1}^N a_{ij} \cdot x_j(t) + \sum_{j=1}^M b_{ij} \cdot u_j(t) + c_i \right)$$

# Computational power

Two models of RNNs are of specific interest:

**RNN[ $\mathbb{Q}$ ]s:** networks with static *rational* weights

**RNN[ $\mathbb{R}$ ]s:** networks with static *real* weights

## Theorem (Siegelmann & Sontag 94, 95)

- ▶ *RNN[ $\mathbb{Q}$ ]s are Turing equivalent.*
- ▶ *RNN[ $\mathbb{R}$ ]s are super-Turing, i.e.  
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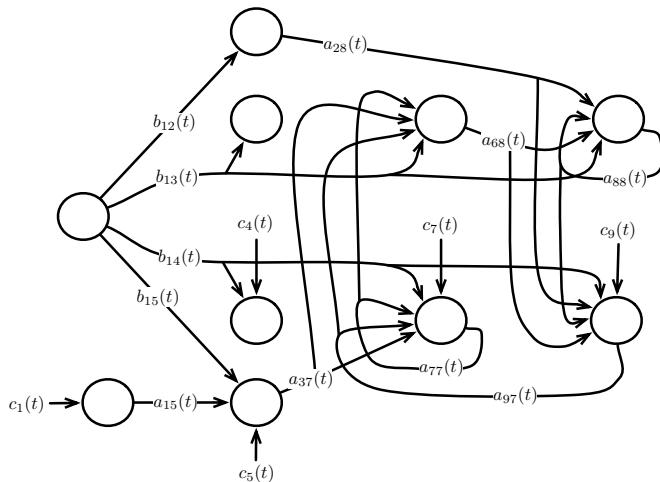
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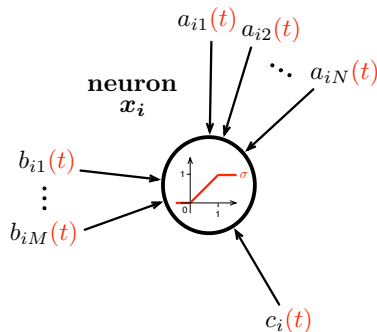
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## Evolving recurrent neural networks



# Dynamics: evolving synaptic weights



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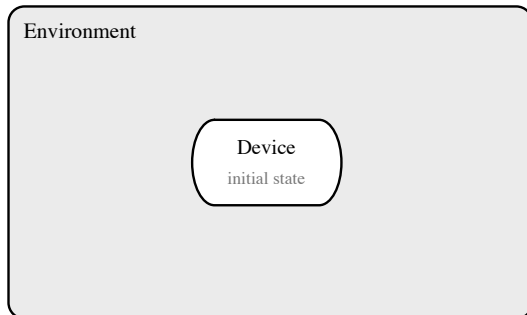
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# Summary of the results

|              | Static  | Evolving                                       |
|--------------|---|--|
| $\mathbb{Q}$ | <b>Turing</b><br>Siegelmann & Sontag 95       | <b>Super-Turing</b><br>Cabessa & Siegelmann 11 |
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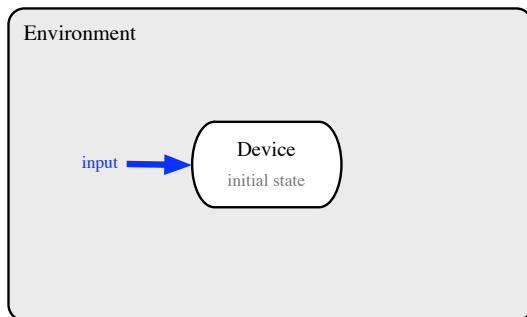
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Closed-box and amnesic...

*"[...] no longer fully corresponds to the current notion of computing in modern systems." (Van Leeuwen & Wiedermann 2008)*

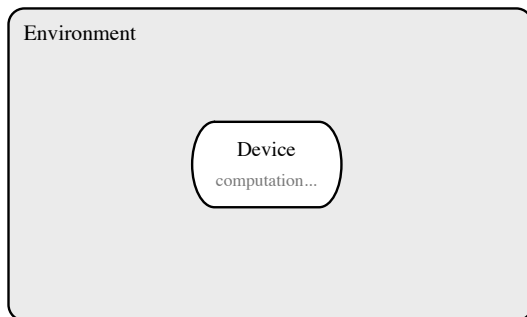
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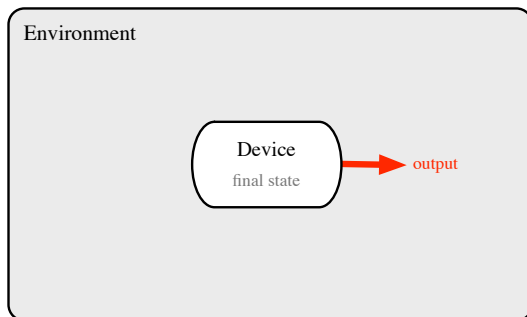
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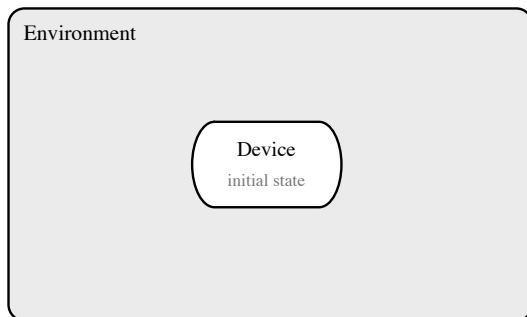
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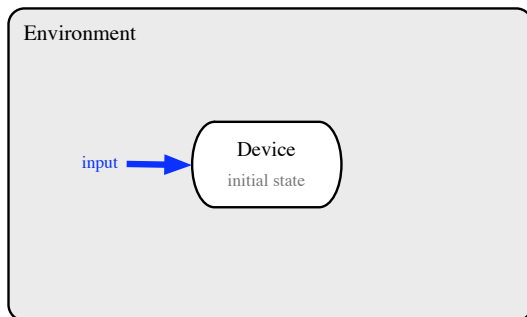


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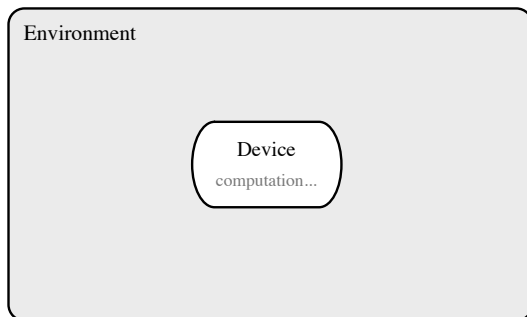
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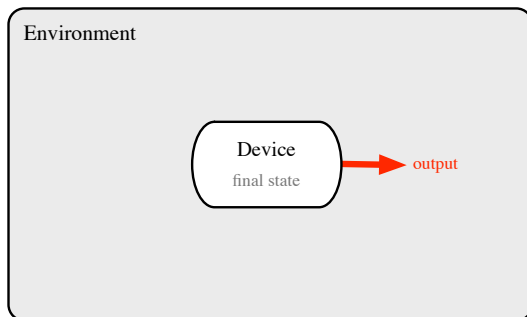
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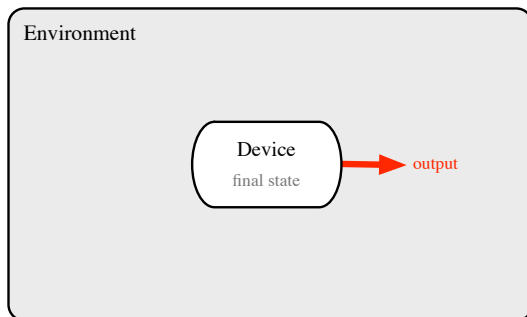
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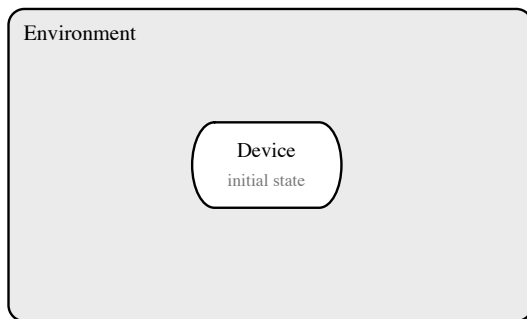
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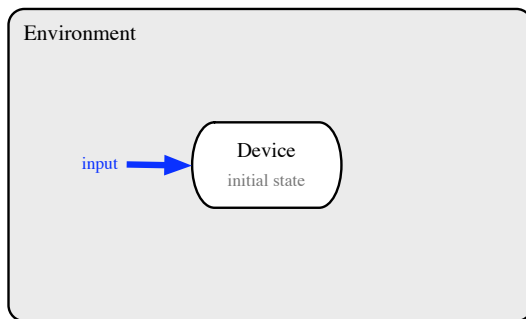
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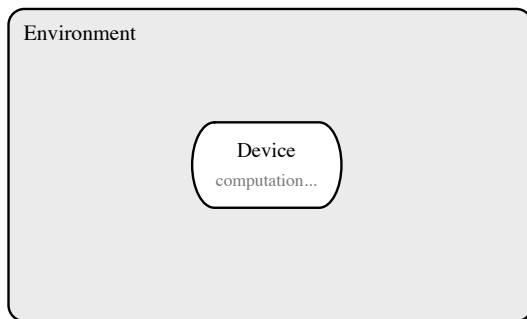
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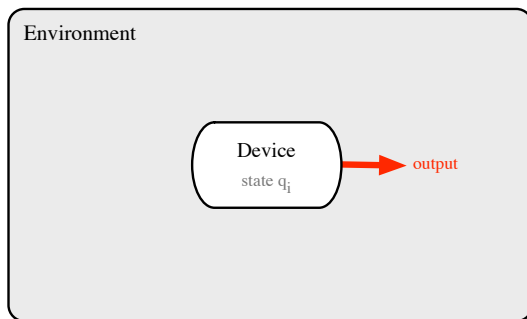
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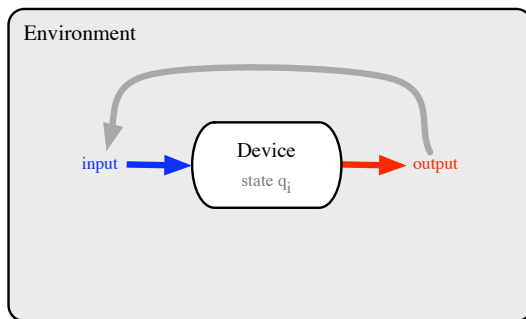
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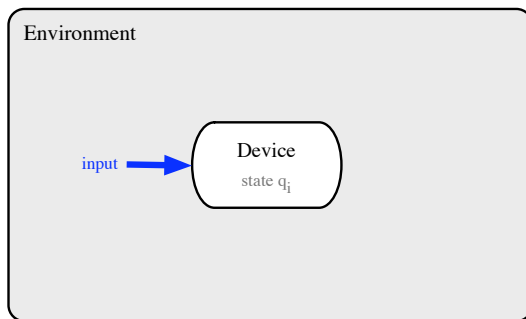


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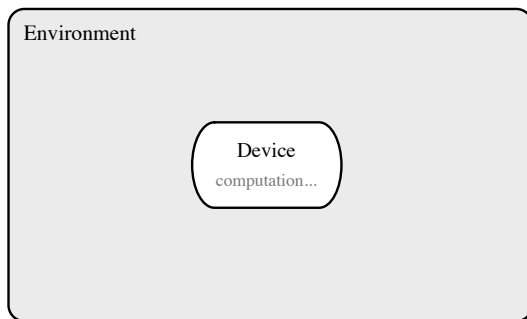
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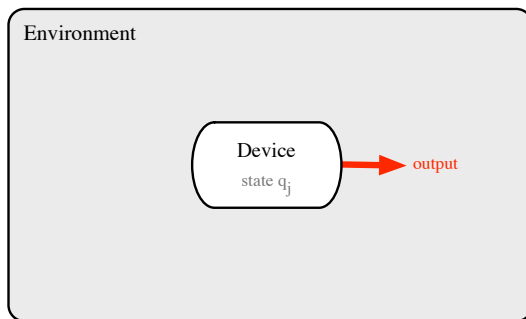
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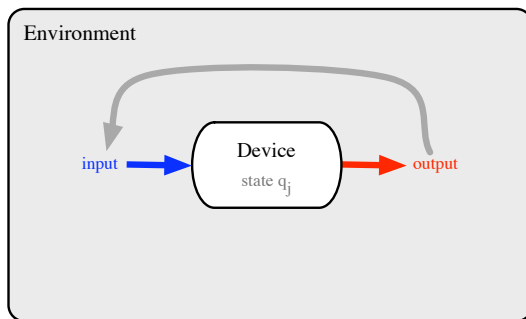
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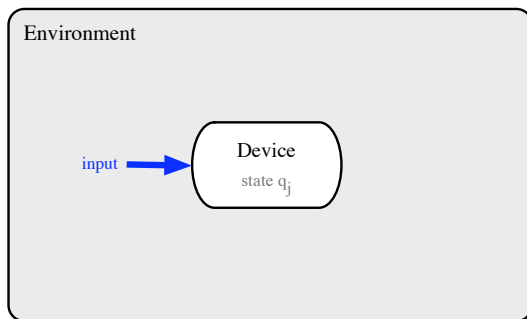
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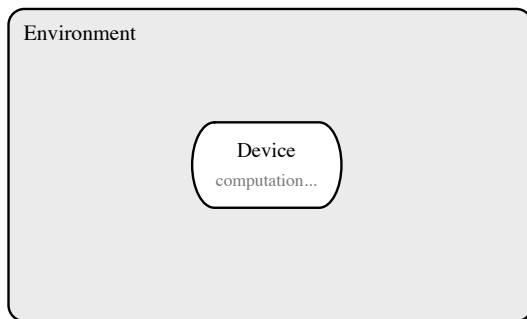
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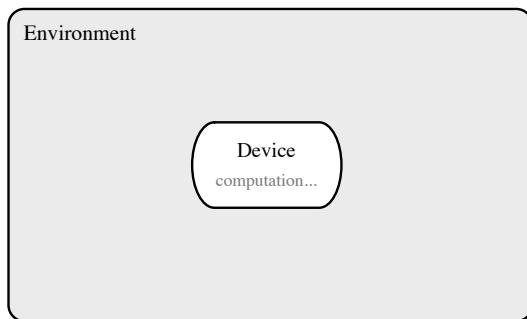
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# Computational power

The results concerning the computational power of RNNs in classical computation generalize to the interactive computational framework.

|              | Static   | Evolving                          |
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# Conclusions

- ▶ Recurrent neural networks provide a natural abstract model of computation beyond the Turing limits.
- ▶ *Architectural Evolution* represents an equivalent alternative to the *power of the continuum* towards the achievement of super-Turing capabilities.
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