

Interactive Evolving Neural Networks are Super-Turing Universal

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Joint work with Alessandro E.P. Villa

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Introduction

- ▶ We follow the so-called *mind-computer analogy* approach to cognitive science.
- ▶ We study the computational capabilities of basic models of recurrent neural networks.
- ▶ We show that recurrent neural networks provide a natural model of computation beyond the Turing limits.

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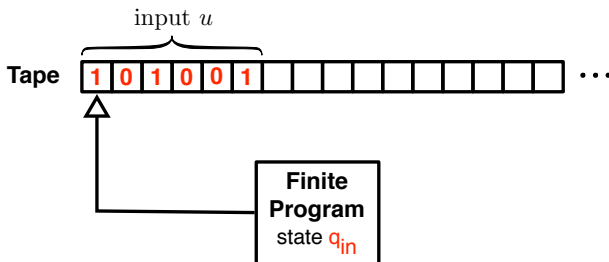
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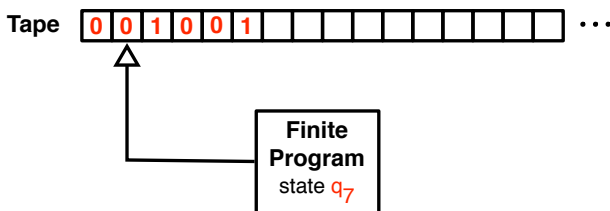
Turing machine

A Turing machine (TM) consists of an infinite tape, a read-write head, and a finite program.



Turing machine

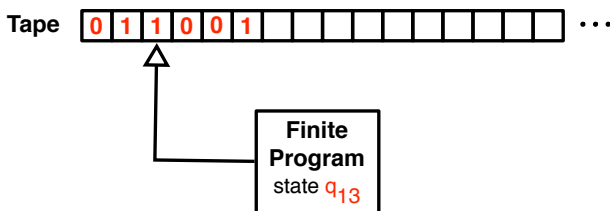
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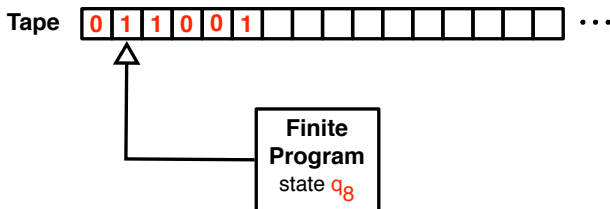
- ▶ input u is *accepted* by \mathcal{M} if $\mathcal{M}(u)$ reaches the state q_{acc}
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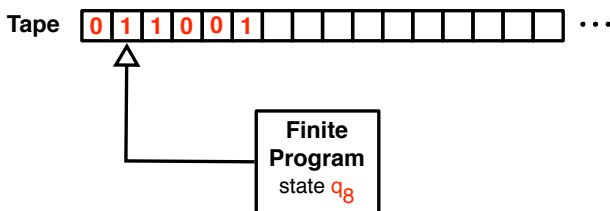


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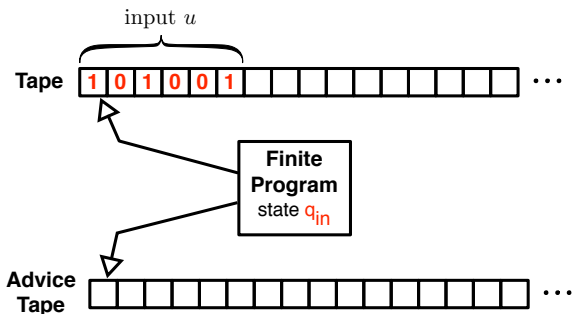


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Turing machine with advice

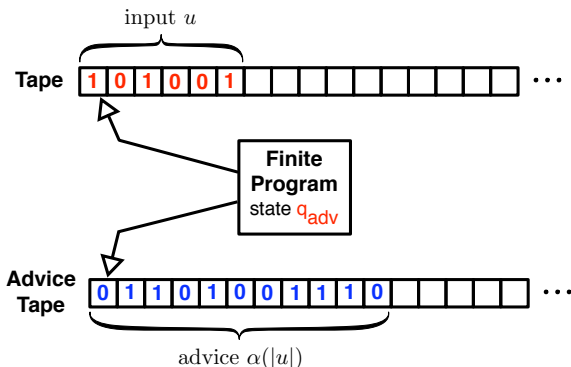
A Turing machine with advice (TM/A) is a Turing machine provided with an additional advice tape and advice function

$$\alpha : \mathbb{N} \longrightarrow \{0, 1\}^*.$$

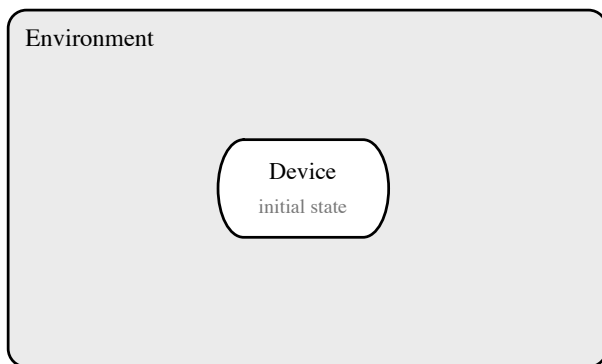


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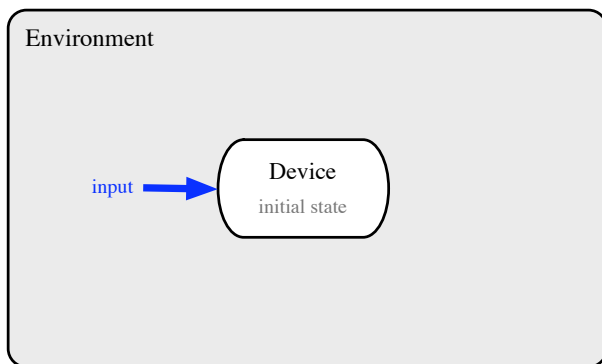


Classical Computation



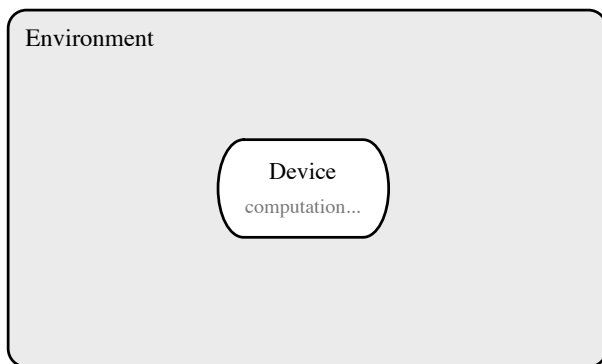
Closed-box and amnesic...

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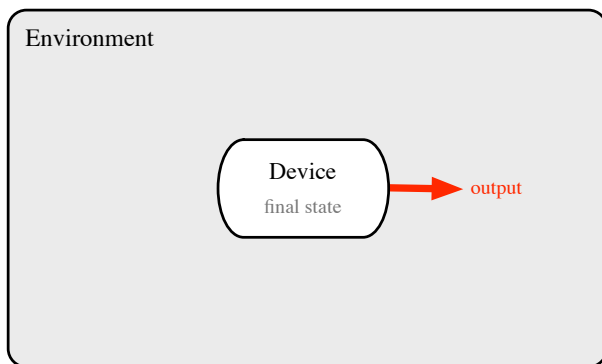
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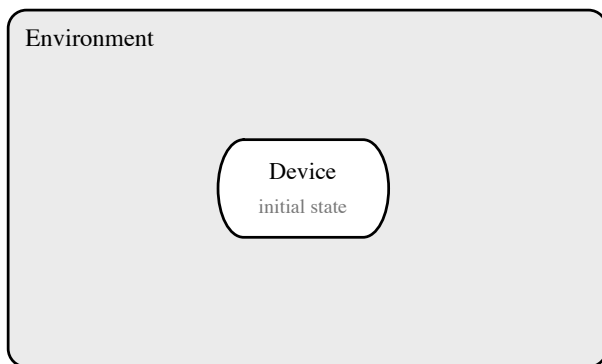
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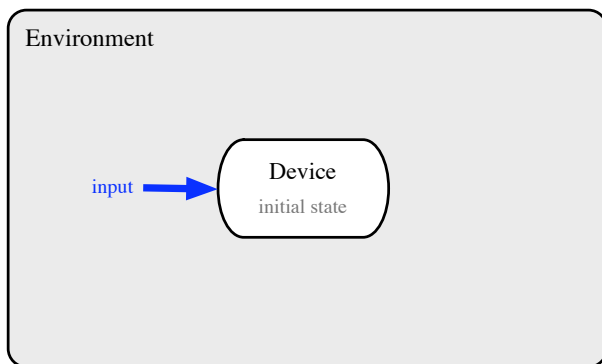
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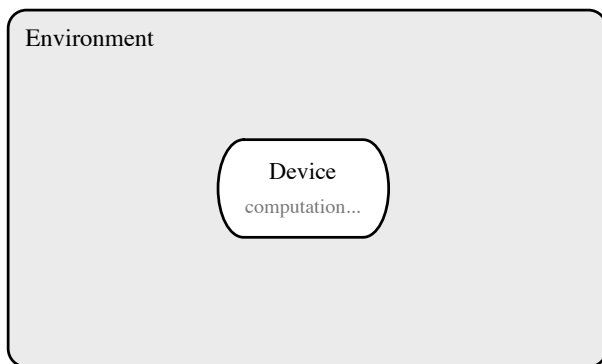
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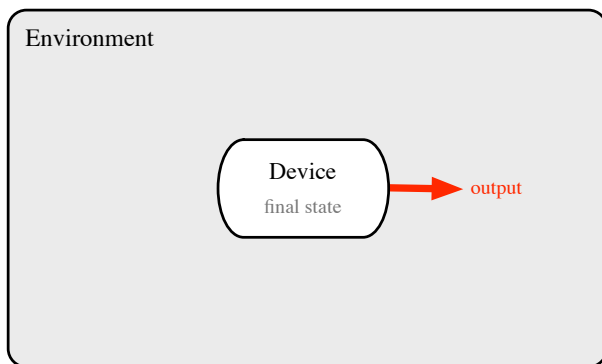
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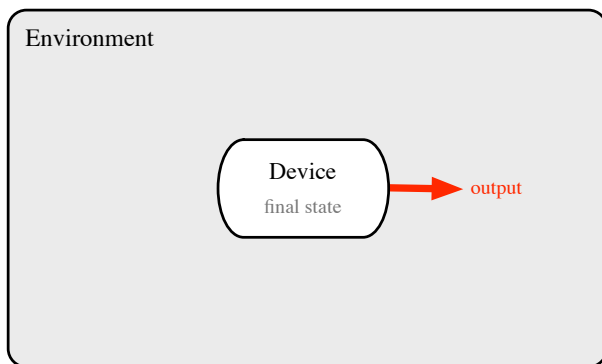
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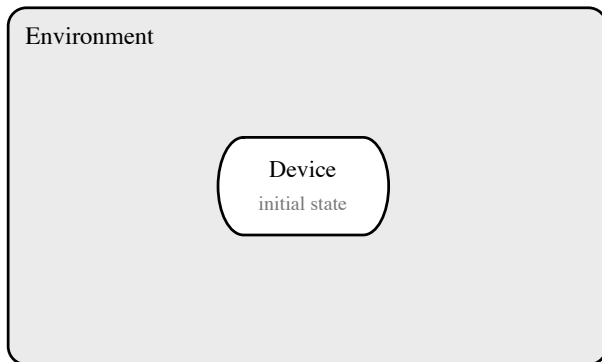
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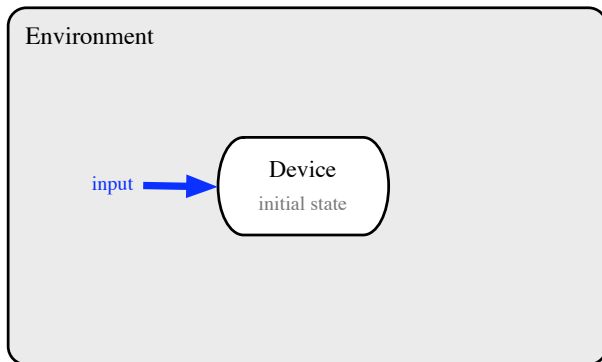
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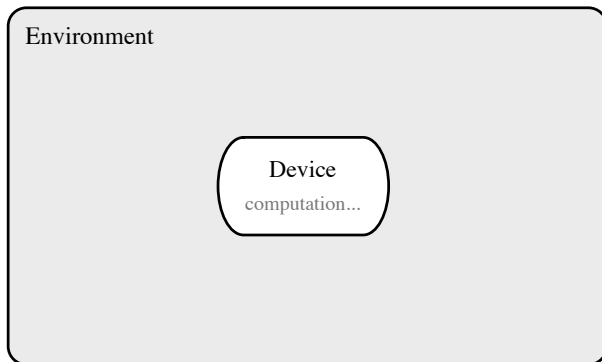
Sequentially interactive and memory active...

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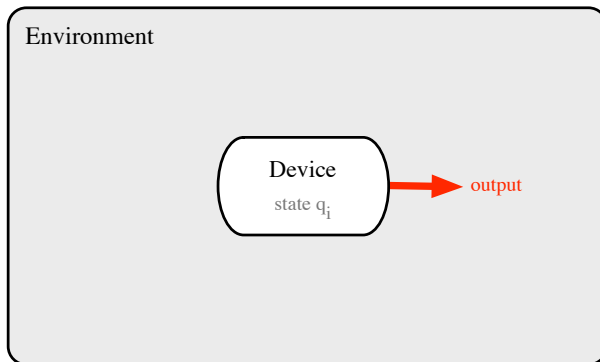
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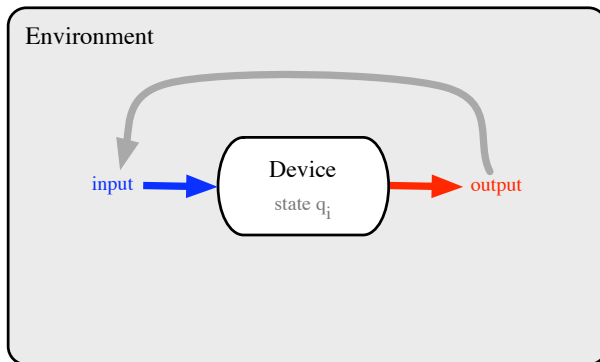
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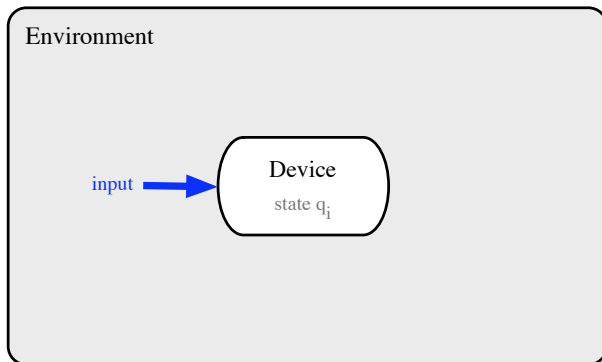
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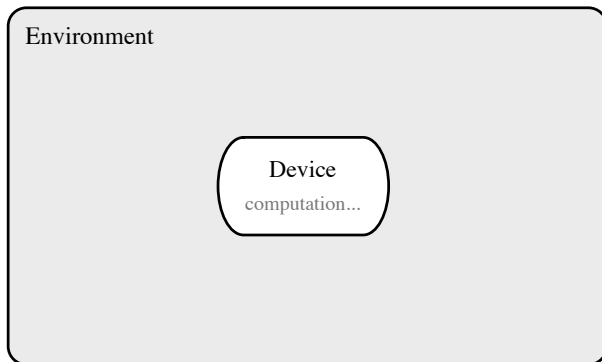
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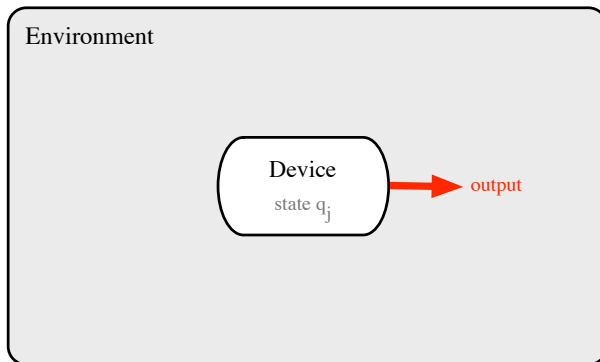
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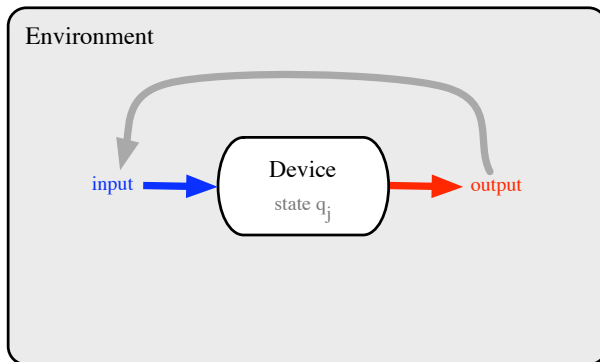
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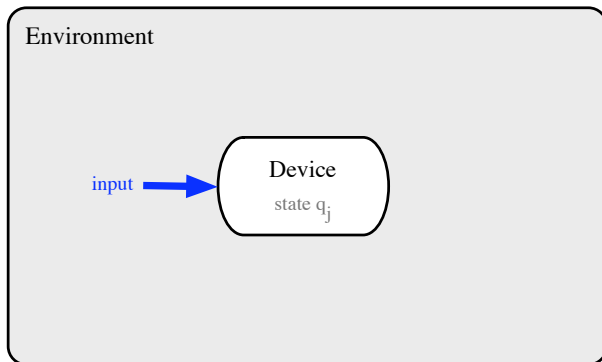
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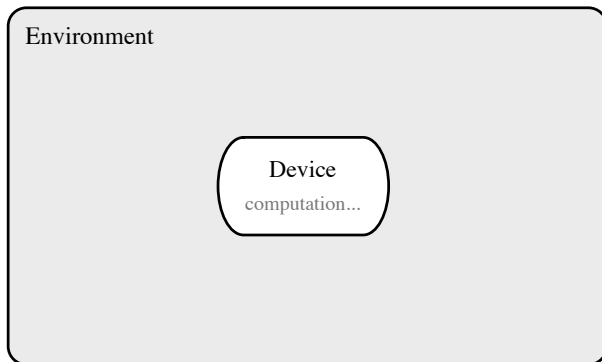
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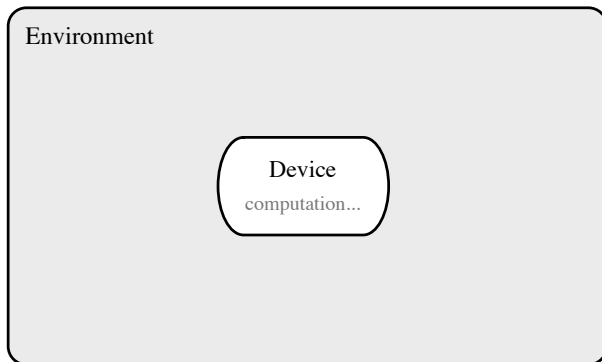
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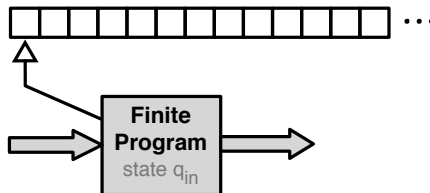
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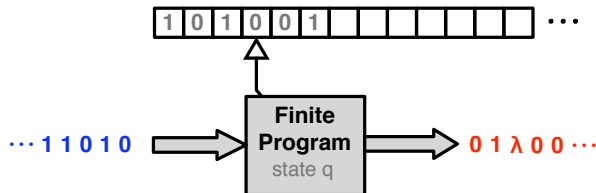


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Interactive Turing machine

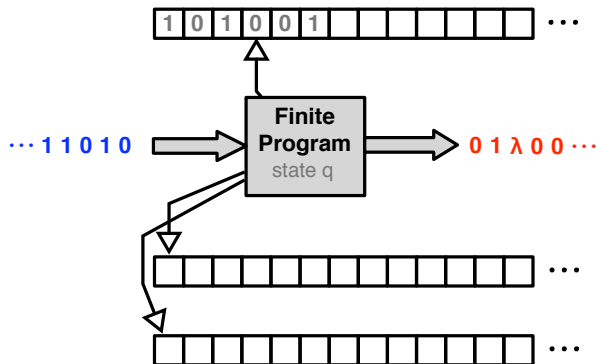


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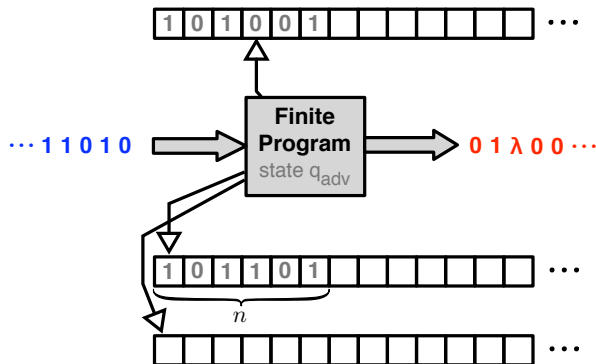
Interactive Turing machine with advice

An Int-TM provided with additional advice input and output tapes and advice function $\alpha : \mathbb{N} \rightarrow \{0, 1\}^*$



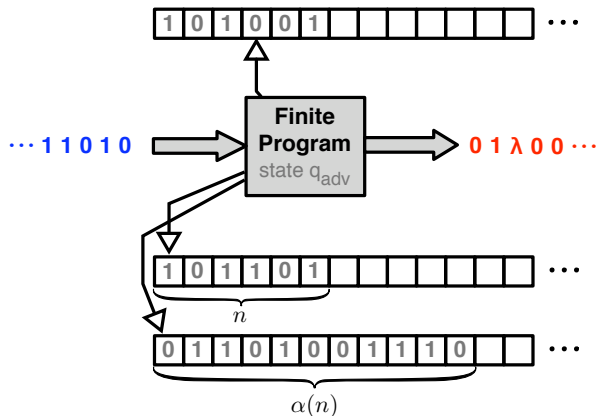
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Lemma

Turing machines with advice are strictly more powerful than Turing machines.

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Interactive Turing machines with advice are strictly more powerful than interactive Turing machines.

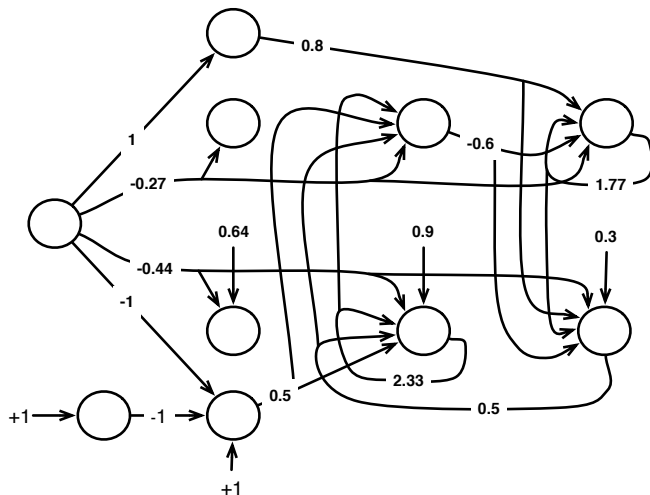
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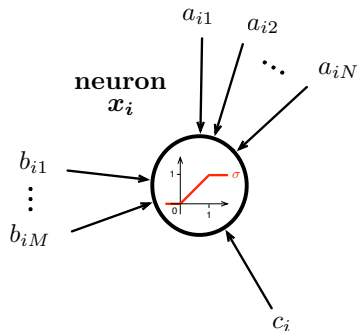
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Recurrent Neural Networks



Dynamics: static synaptic weights

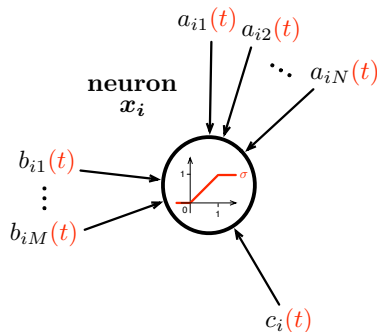


$$x_i(t+1) = \sigma \left(\sum_{j=1}^N a_{ij} \cdot x_j(t) + \sum_{j=1}^M b_{ij} \cdot u_j(t) + c_i \right)$$

Results

	Static Architecture
\mathbb{Q}	Turing Siegelmann & Sontag 95 (classical comp.) Cabessa & Siegelmann 12 (interactive comp.)
\mathbb{R}	Super-Turing Siegelmann & Sontag 94 (classical comp.) Cabessa & Siegelmann 12 (interactive comp.) Cabessa & Villa 12 (interactive comp.)

Dynamics: evolving synaptic weights

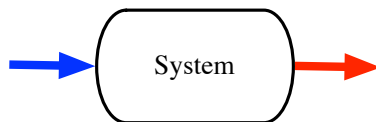


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	Evolving Architecture
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Interactive Scenario

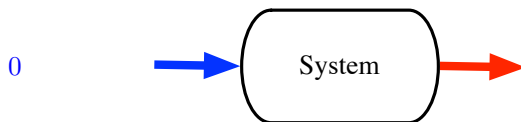


In this way, any deterministic interactive system S realises an ω -translation $\varphi_S : \{0, 1\}^\omega \rightarrow \{0, 1\}^{\leq \omega}$.

Definition

An ω -translation ψ is said to be *interactively computable* if there exists a deterministic interactive system S (of any kind) such that $\varphi_S = \psi$.

Interactive Scenario

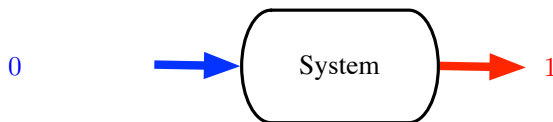


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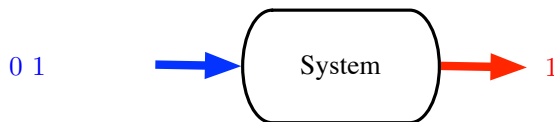


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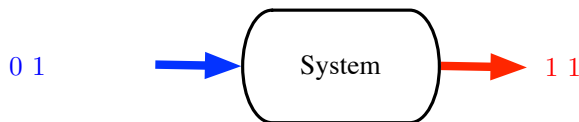


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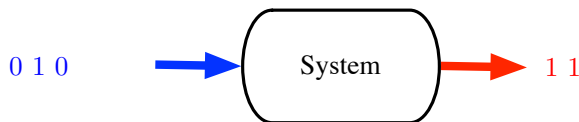


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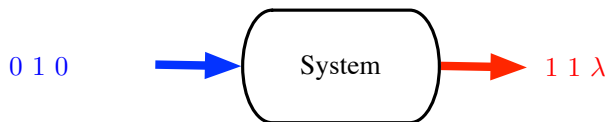


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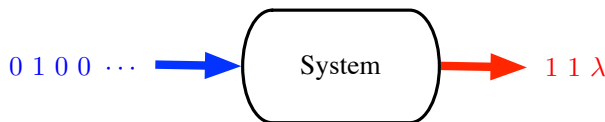


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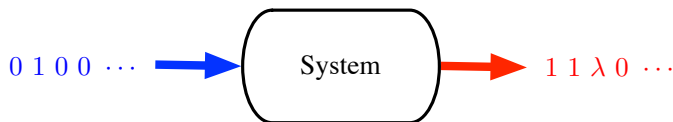


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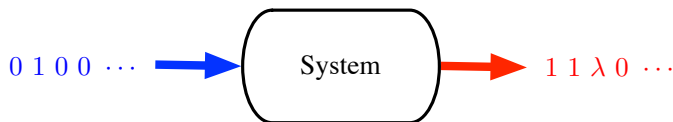


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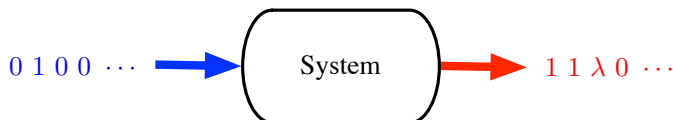


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Theorem

Interactive Evolving RNNs are Super-Turing universal.

More precisely, let ψ be some ω -translation. The following conditions are equivalent:

- 1. ψ is interactively computable (by any det. int. system)*
- 2. ψ is computable by some ω -E-RNN*
- 3. ψ is computable by some ω -E-RNN with ω -input*
- 4. ψ is computable by some ω -E-RNN?*
- 5. ψ is computable by some ω -TSA*

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3. ψ is realisable by some Int-Ev-RNN[\mathbb{Q}]
4. ψ is realisable by some Int-St-RNN[\mathbb{R}]
5. ψ is realisable by some Int-TM/A
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Summary

	Static	Evolving
\mathbb{Q}	Turing	Super-Turing universal (interactive case)
\mathbb{R}	Super-Turing universal (interactive case)	Super-Turing universal (interactive case)

Conclusions

- ▶ The results are theoretical: any implementation of an evolving RNN will never be super-Turing.
- ▶ Evolving-RNNs provide a natural abstract computational model beyond the Turing limits.
- ▶ *Architectural Evolution* is an alternative way to the *power of the continuum* to achieve super-Turing capabilities.
- ▶ The results support the idea that *architectural evolution* might play a crucial role in the computational capabilities of biological neural networks.
- ▶ Future work: study the computational power of more biologically oriented neural models involved in more bio-inspired computational frameworks.
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